

Making the most out of noisy quantum computers: Strategies for circuit design and error mitigation

Stefan Kühn



EUROCC
MAY 2021

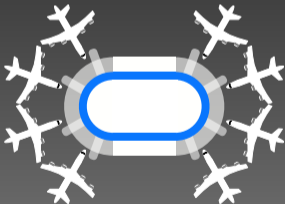
Motivation

Problems for which quantum computers might be advantageous

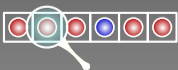
- Factoring

$$70747 = 263 \times 269$$

- Optimization

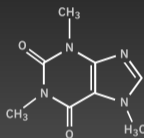


- Searching databases

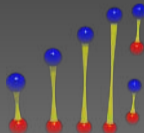


- Quantum simulation

- ▶ Quantum chemistry



- ▶ Particle physics

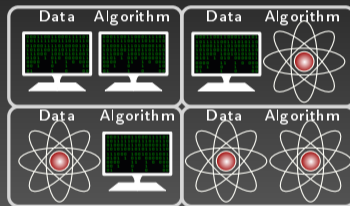


- ▶ Cosmology

- ▶ Material science

- ▶ ...

- Machine learning



- Cryptography

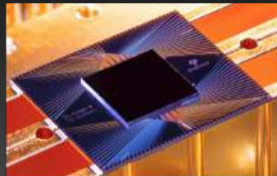


- ...

Motivation

On the verge of the NISQ era

- Noisy Intermediate-Scale Quantum (NISQ) technology (50-100 qubits) is already available
- First small/scale devices are commercially available
- Noise significantly limits the circuit depths that can be executed reliably
 - ▶ Not fault tolerant
 - ▶ No quantum error correction



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Article

Quantum supremacy using a programmable superconducting processor

RESEARCH

<https://doi.org/10.1038/s41585-019-0101-1>

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QUANTUM COMPUTING

Quantum computational advantage using photons

Han-Sen Zhong^{1,2*}, Hui Wang^{1,2,*}, Yu-Hao Deng^{1,2,*}, Ming-Cheng Chen^{1,2,*}, Li-Chao Peng^{1,2}, Yi-Han Luo^{1,2}, Jian Qin^{1,2}, Dian Wu^{1,2}, Xing Ding^{1,2}, Yi Hu^{1,2}, Peng Hu¹, Xiao-Yan Yang¹, Wei-Jun Zhang¹, Hao Li¹, Yuxuan Li¹, Xiao Jiang^{1,2}, Lin Gan¹, Guangwen Yang¹, Liang You¹, Zhen Wang¹, Li Li^{1,2}, Nai-Le Liu^{1,2}, Chao-Yang Lu^{1,2}, Jian-Wei Pan^{1,2}

Quantum computers promise to perform certain tasks that are believed to be intractable to classical computers. Boson sampling is such a task and is considered a strong candidate to demonstrate the quantum computational advantage. We performed Gaussian boson sampling by sending 50 indistinguishable single-mode squeezed states into a 100-mode ultralow-loss interferometer with full connectivity and random matrix—the whole optical setup is phase-locked—and sampling the output using 100 high-efficiency single-photon detectors. The obtained samples were validated against plausible hypotheses exploiting thermal states, distinguishable photons, and uniform distribution. The photonic quantum computer, *Aurang*, generates up to 76 output photon clicks, which yields an output state-space dimension of 20^{76} and a sampling rate that is faster than using the state-of-the-art simulation strategy and supercomputers by a factor of $\sim 10^{14}$.

F. Arute et al., Nature 574, 5050 (2018)
H.-S. Zhong et al., Science 370, 1460 (2020)

Outline

- 1 Motivation
- 2 The Circuit Model of Quantum Computing
- 3 Dimensional Expressivity Analysis
- 4 Measurement Error Mitigation
- 5 Summary & Outlook

The Circuit Model of Quantum Computing

Quantum bit

- Qubit: two-dimensional quantum system
- Hilbert space \mathcal{H} with basis $\{|0\rangle, |1\rangle\}$, called the **computational basis**
- Contrary to classical bits, it can be in a **superposition**

$$|\psi\rangle = \alpha |0\rangle + \beta |1\rangle, \quad \alpha, \beta \in \mathbb{C}, \quad |\alpha|^2 + |\beta|^2 = 1$$

The Circuit Model of Quantum Computing

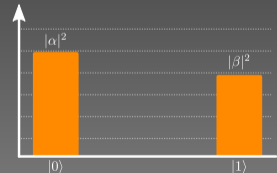
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- Measuring the qubit in the computational basis collapses the state onto one of the basis states
- Probabilities of measuring the two outcomes 0 and 1

$$\begin{aligned} p(0) &= |\alpha|^2, & |\psi\rangle &= |0\rangle \\ p(1) &= |\beta|^2, & |\psi\rangle &= |1\rangle \end{aligned}$$



- Unlike for classical systems **measurement changes the quantum system**

The Circuit Model of Quantum Computing

Multiple quantum bits

- N qubits: Hilbert space is the tensor product $\underbrace{\mathcal{H} \otimes \cdots \otimes \mathcal{H}}_{N \text{ times}}$
- Most general state in the computational basis

$$|\psi\rangle = \sum_{i_1, \dots, i_N=0}^1 c_{i_1 \dots i_N} |i_1\rangle \otimes \cdots \otimes |i_N\rangle$$

- Multiple qubits can be **entangled**

The Circuit Model of Quantum Computing

Entanglement of bipartite systems

- Consider bipartite systems $\mathcal{H}_A \otimes \mathcal{H}_B$
- A quantum state that can be factored as a tensor product of states of its local constituents is called **separable state** or **product state**

$$|\psi\rangle = |\psi_A\rangle \otimes |\psi_B\rangle$$

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Example

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- $|\Phi^+\rangle = \frac{1}{\sqrt{2}} (|0\rangle \otimes |0\rangle + |1\rangle \otimes |1\rangle)$
 \Rightarrow entangled state (Bell state)

The Circuit Model of Quantum Computing

Entanglement of bipartite systems

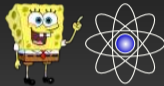
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The Circuit Model of Quantum Computing

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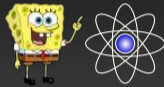
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- Bob can measure his qubit

$$p_{\text{Bob}}(0) = \frac{1}{2},$$

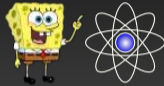
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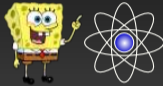
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$$p_{\text{Bob}}(0) = \frac{1}{2},$$

$$|\psi\rangle = |0\rangle \otimes |0\rangle,$$

$$p_{\text{Alice}}(0) = 1$$

$$p_{\text{Bob}}(1) = \frac{1}{2},$$

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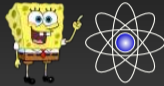
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- If Alice measures after Bob she obtains the same result with certainty
- ⇒ The measurement outcomes are perfectly correlated

The Circuit Model of Quantum Computing

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⇒ The measurement outcomes are perfectly correlated

- Choice of measurement at one location affects the other qubit

⇒ “Spooky action at a distance”

The Circuit Model of Quantum Computing

Quantum gates

- Quantum mechanics is reversible, $|\psi\rangle$ undergoes unitary evolution under some (time-dependent) Hamiltonian $\mathcal{H}(t)$

$$|\psi(t)\rangle = \underbrace{T \exp\left(-i \int_0^t ds \mathcal{H}(s)\right)}_{\text{unitary matrix } U} |\psi_0\rangle$$

- Quantum gates are represented by **unitary matrices**

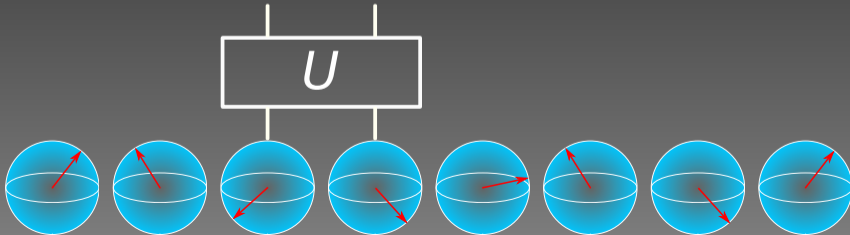
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- **Quantum gates** are represented by **unitary matrices**
- Typically gates only act on a few qubits in a nontrivial way



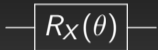
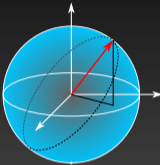

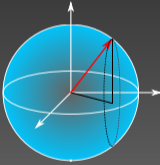

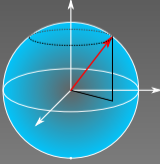
The Circuit Model of Quantum Computing

Common single-qubit quantum gates

Hadamard		$H = \begin{pmatrix} \frac{1}{\sqrt{2}} & \frac{1}{\sqrt{2}} \\ \frac{1}{\sqrt{2}} & -\frac{1}{\sqrt{2}} \end{pmatrix}$	$\begin{aligned} 0\rangle &\rightarrow \frac{1}{\sqrt{2}} (0\rangle + 1\rangle) \\ 1\rangle &\rightarrow \frac{1}{\sqrt{2}} (0\rangle - 1\rangle) \end{aligned}$
X		$X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}$	$\begin{aligned} 0\rangle &\rightarrow 1\rangle \\ 1\rangle &\rightarrow 0\rangle \end{aligned}$
Y		$Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}$	$\begin{aligned} 0\rangle &\rightarrow -i 1\rangle \\ 1\rangle &\rightarrow i 0\rangle \end{aligned}$
Z		$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$	$\begin{aligned} 0\rangle &\rightarrow 0\rangle \\ 1\rangle &\rightarrow - 1\rangle \end{aligned}$

The Circuit Model of Quantum Computing

Common single-qubit rotations

$R_X(\theta)$		$R_X(\theta) = \exp\left(-i\frac{\theta}{2}X\right)$	
$R_Y(\theta)$		$R_Y(\theta) = \exp\left(-i\frac{\theta}{2}Y\right)$	
$R_Z(\theta)$		$R_Z(\theta) = \exp\left(-i\frac{\theta}{2}Z\right)$	

The Circuit Model of Quantum Computing

Common multi-qubit quantum gates

CNOT



$$\text{CNOT} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{pmatrix}$$

$$|0\rangle \otimes |0\rangle \rightarrow |0\rangle \otimes |0\rangle$$




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The Circuit Model of Quantum Computing

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SWAP gate		$\text{SWAP} = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 \\ 0 & 0 & 0 & 1 \end{pmatrix}$	$ 0\rangle \otimes 0\rangle \rightarrow 0\rangle \otimes 0\rangle$ $ 0\rangle \otimes 1\rangle \rightarrow 1\rangle \otimes 0\rangle$ $ 1\rangle \otimes 0\rangle \rightarrow 0\rangle \otimes 1\rangle$ $ 1\rangle \otimes 1\rangle \rightarrow 1\rangle \otimes 1\rangle$
Toffoli gate		$\text{Toffoli} = \begin{pmatrix} \mathbb{1}_{6 \times 6} & 0 & 0 \\ 0 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}$	CCNOT

The Circuit Model of Quantum Computing

Quantum gates

- The reversible classical gates can be implemented on a quantum computer
⇒ We can replicate classical computation

The Circuit Model of Quantum Computing

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- The Hadamard gate can **create superpositions** out of a single basis state

$$|0\rangle \text{ --- } \boxed{H} \text{ --- } |+\rangle \qquad |0\rangle \rightarrow |+\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle)$$

- The CNOT gate can **create entanglement**

$$\begin{array}{c} |\psi_1\rangle \text{ --- } \bullet \\ |\psi_2\rangle \text{ --- } \oplus \end{array} \qquad |\psi_1\rangle \otimes |\psi_2\rangle \rightarrow |\phi_{12}\rangle \neq |\phi_1\rangle \otimes |\phi_2\rangle$$

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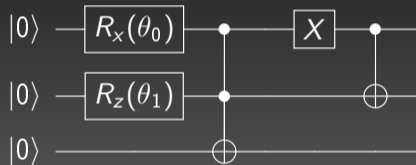
- Since quantum mechanics is linear, we can **apply gates to superpositions of basis states**

$$\begin{aligned} \text{CNOT}(\alpha|0\rangle \otimes |0\rangle + \beta|0\rangle \otimes |1\rangle + \gamma|1\rangle \otimes |0\rangle + \delta|1\rangle \otimes |1\rangle) \\ = \alpha|0\rangle \otimes |0\rangle + \beta|0\rangle \otimes |1\rangle + \gamma|1\rangle \otimes |1\rangle + \delta|1\rangle \otimes |0\rangle \end{aligned}$$

The Circuit Model of Quantum Computing

Quantum circuits

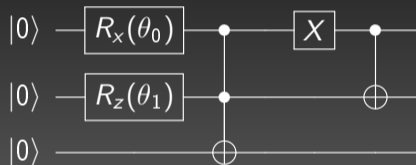
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The Circuit Model of Quantum Computing

Quantum circuits

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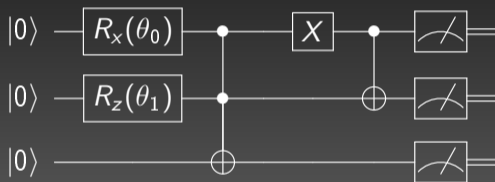


- Depth** of a circuit: maximum length of a directed path from the input to the output

The Circuit Model of Quantum Computing

Quantum circuits

- Combining multiple gates we can build **quantum circuits**



- Depth** of a circuit: maximum length of a directed path from the input to the output
- Extracting information: final measurement of the qubits (usually in the computational basis)

The Circuit Model of Quantum Computing

Why is Quantum Computing more powerful?

- The Hilbert space of N qubits is the tensor product $\underbrace{\mathcal{H} \otimes \cdots \otimes \mathcal{H}}_{N \text{ times}}$

\Rightarrow Dimension 2^N , **number of basis states grows exponentially**

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\Rightarrow “Quantum parallelism”

The Circuit Model of Quantum Computing

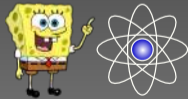
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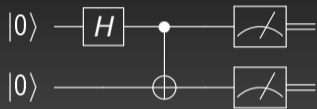


\Rightarrow **Correlations** that have no classical analog

The Circuit Model of Quantum Computing

Example

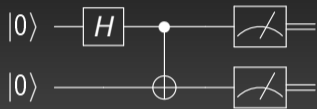
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The Circuit Model of Quantum Computing

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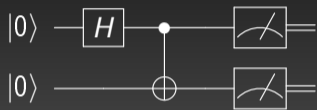


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The Circuit Model of Quantum Computing

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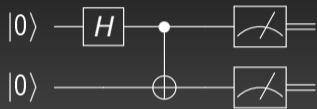
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The Circuit Model of Quantum Computing

Example

- Simple circuit preparing an entangled state (Bell state)



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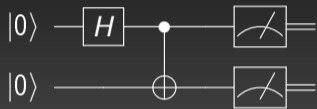
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The Circuit Model of Quantum Computing

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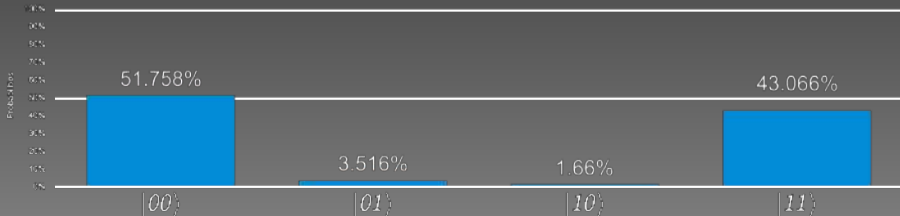


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- Results on actual quantum hardware (ibmq_vigo)



The Circuit Model of Quantum Computing

Where do we stand?

Current NISQ devices

- Small or intermediate scale
- Considerable amount of noise
- Only shallow circuits can be executed faithfully
- Quantum advantage demonstrated

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How can we utilize existing quantum hardware in a beneficial way?

The Circuit Model of Quantum Computing

Hybrid quantum-classical algorithms

- Combine classical and quantum devices
- Rely on classical computing where possible
- Use the quantum device as a coprocessor
 - ▶ Tackle the classically hard/intractable part of the problem
 - ▶ Feed the classical data obtained from a measurement back to the classical computer



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Even modest quantum hardware can yield advantages

The Circuit Model of Quantum Computing

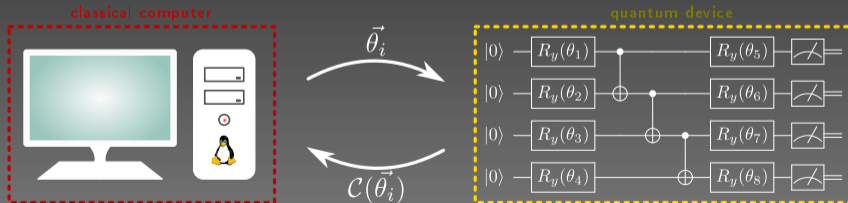
Variational Quantum Eigensolver

- Algorithm to find ground states of quantum Hamiltonians \mathcal{H}
- Define a cost function $\mathcal{C}(\vec{\theta}) = \langle \psi(\vec{\theta}) | \mathcal{H} | \psi(\vec{\theta}) \rangle$, $\vec{\theta} = \mathbb{R}^n$

The Circuit Model of Quantum Computing

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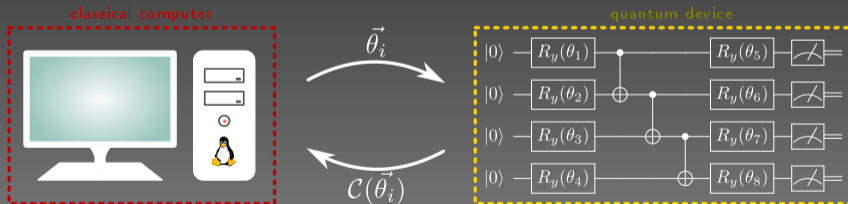
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The Circuit Model of Quantum Computing

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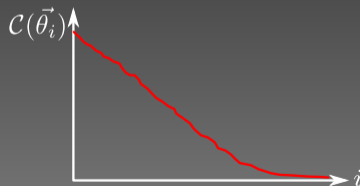
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The Circuit Model of Quantum Computing

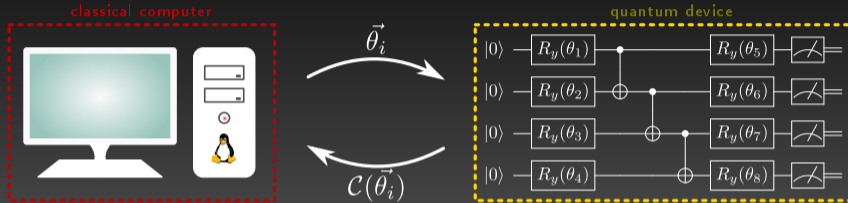
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The Circuit Model of Quantum Computing

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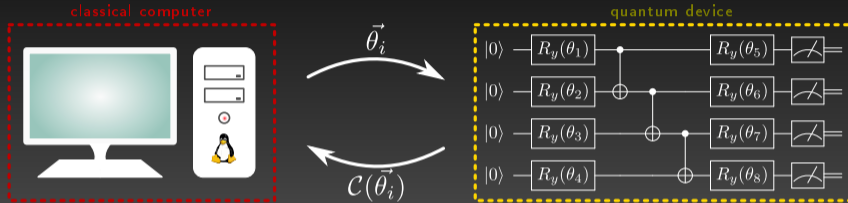


Advantages

- Flexible ansatz design
- Hamiltonian exists only as a measurement
- Partially resilient to systematic errors

The Circuit Model of Quantum Computing

Variational Quantum Algorithms



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Challenges

- How to choose an expressive ansatz?
- How to avoid redundant parameters?
- How to deal with effects of noise?

3.

- 1 Motivation
- 2 The Circuit Model of Quantum Computing
- 3 Dimensional Expressivity Analysis
- 4 Measurement Error Mitigation
- 5 Summary & Outlook

Dimensional Expressivity Analysis

Number of parameters in the ansatz circuit should be

- **large** for solutions to be reachable
- **large** in order not to introduce artificial local optima
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Optimal circuit for VQE

- **maximally expressive**: be able to generate all (physically relevant) states
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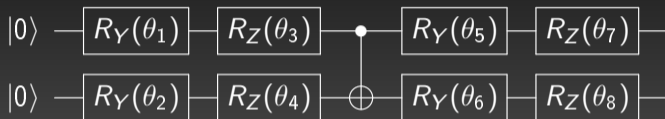
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Can we develop a mathematical framework to determine if a circuit is both minimal and maximally expressive?

Dimensional Expressivity Analysis

Dimensional Expressivity Analysis

- Parametric quantum circuit with parameters $\vec{\theta} \in P \subseteq \mathbb{R}^n$ generating $|C(\vec{\theta})\rangle$



- Treat the parametric circuit as a map that maps the input parameters to the state space of the quantum device

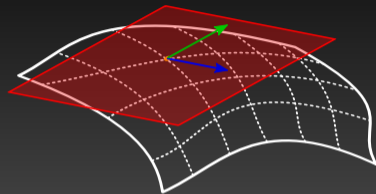
$$C : \vec{\theta} \mapsto |C(\vec{\theta})\rangle = R_Z(\theta_8) \dots R_Y(\theta_1) |0\rangle \otimes |0\rangle$$

- Parameter space P : real manifold
- Image of C : **circuit manifold** \mathcal{M}
- Which parameters are necessary to generate the circuit manifold \mathcal{M} ?

Dimensional Expressivity Analysis

Dimensional Expressivity Analysis

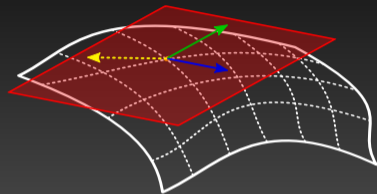
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Dimensional Expressivity Analysis

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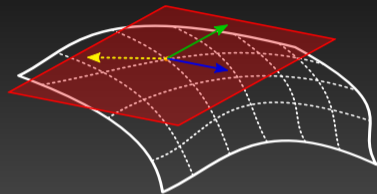
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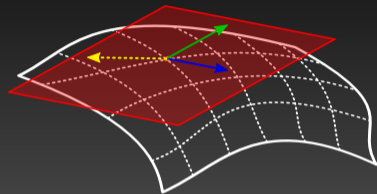
Iterative procedure to identify redundant parameters

- θ_1 is never redundant as long as the corresponding parametric gate is nontrivial
- Check whether $|\partial_{k+1} C(\vec{\theta})\rangle$ is a linear combination of $|\partial_1 C(\vec{\theta})\rangle, \dots, |\partial_k C(\vec{\theta})\rangle$
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- Remove redundant parameters
 - ▶ Parameter removal implies setting the parameter to a constant value
 - ▶ Rotation gates (e.g. $\exp(-\frac{i}{2}\vartheta X)$): choose the parameter $\vartheta = 0$ to achieve an $\mathbb{1}$

Dimensional Expressivity Analysis

Checking for parameter independence

- θ_1 is never redundant as long as corresponding parametric gate is nontrivial
- For θ_k , $k = 2, \dots, n$ repeat the following steps
 - ▶ Since P is a real manifold, we have to consider the real Jacobian

$$J_k = \begin{pmatrix} \Re|\partial_1 C\rangle & \dots & \Re|\partial_k C\rangle \\ \Im|\partial_1 C\rangle & \dots & \Im|\partial_k C\rangle \end{pmatrix}$$

- ▶ If the matrix J_k has full rank then θ_k is independent
- Instead of checking the rank of J_k one can also compute the rank of $S_k = J_k^T J_k$

Dimensional Expressivity Analysis

Dimensional Expressivity Analysis

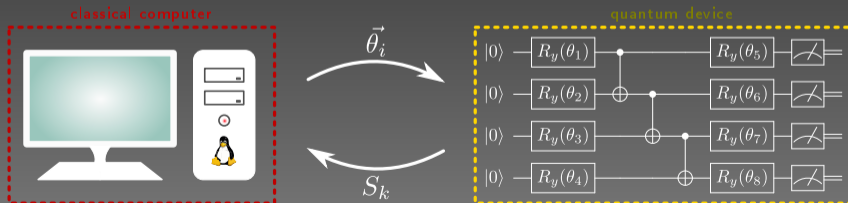
- **Memory requirements on a classical computer: exponential in the number of qubits N because J_k has dimensions $2^{N+1} \times k$**

Dimensional Expressivity Analysis

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Can we use a hybrid-quantum classical approach for the Dimensional Expressivity Analysis?



Dimensional Expressivity Analysis

Hybrid Quantum-Classical Dimensional Expressivity Analysis

- Since the first parameter is always nontrivial $S_1 = \frac{1}{4}$
- For $k \geq 2$ the $k \times k$ matrices $S_k = J_k^T J_k$ can be cast into the form

$$S_k = \begin{pmatrix} S_{k-1} & A_k \\ A_k^T & \frac{1}{4} \end{pmatrix} \quad \text{with} \quad A_k = \begin{pmatrix} \Re \langle \partial_1 C(\vec{\theta}) | \partial_k C(\vec{\theta}) \rangle \\ \vdots \\ \Re \langle \partial_{k-1} C(\vec{\theta}) | \partial_k C(\vec{\theta}) \rangle \end{pmatrix}$$

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- For $R_G(\vartheta) = \exp(-\frac{i}{2}\vartheta G)$ where G is a gate, the derivative is essentially a circuit

$$|R_G\rangle = |0\rangle \text{ --- } \boxed{R_G(\vartheta)} \text{ ---} \quad \Rightarrow \quad 2i |\partial_\theta R_G\rangle = |0\rangle \text{ --- } \boxed{R_G(\vartheta)} \text{ --- } \boxed{G} \text{ ---}$$

- Up to an imaginary factor $|\partial_j C(\vec{\theta})\rangle$ can be prepared on a quantum device

Dimensional Expressivity Analysis

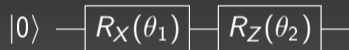
Hybrid quantum-classical Dimensional Expressivity Analysis

- If we can efficiently obtain $\Re\langle\partial_j C(\vec{\theta})|\partial_k C(\vec{\theta})\rangle$ on the quantum device, we can carry out dimensional expressivity analysis efficiently

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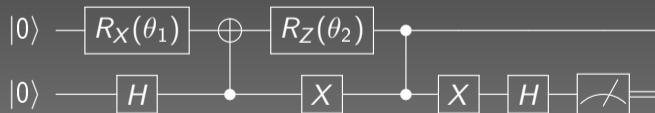
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- Circuit for obtaining $\Re\langle\partial_1 C(\vec{\theta})|\partial_2 C(\vec{\theta})\rangle$



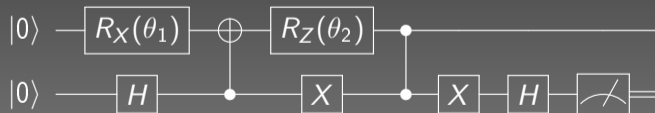
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- Circuit for obtaining $\Re\langle\partial_1 C(\vec{\theta})|\partial_2 C(\vec{\theta})\rangle$



- Real part of the overlap is proportional to the probability for the ancilla being in $|0\rangle$

S. Lloyd, M. Mohseni, P. Rebentrost, arXiv:1307.0411 (2013)

L. Zhao, Z. Zhao, P. Rebentrost, J. Fitzsimons, arXiv:1902.10394 (2019)

Lena Funcke, Tobias Hartung, Karl Jansen, SK, Paolo Stornati, Quantum 5, 422 (2021)

Dimensional Expressivity Analysis

Results for a single qubit on quantum hardware

- Circuit we examine

$$C(\theta_4, \theta_3, \theta_2, \theta_1) = \\ R_Y(\theta_4)R_Z(\theta_3)R_X(\theta_2)R_Z(\theta_1) |0\rangle$$

- Number of independent parameters: 3
- $R_Y(\theta_4)$ is redundant

Dimensional Expressivity Analysis

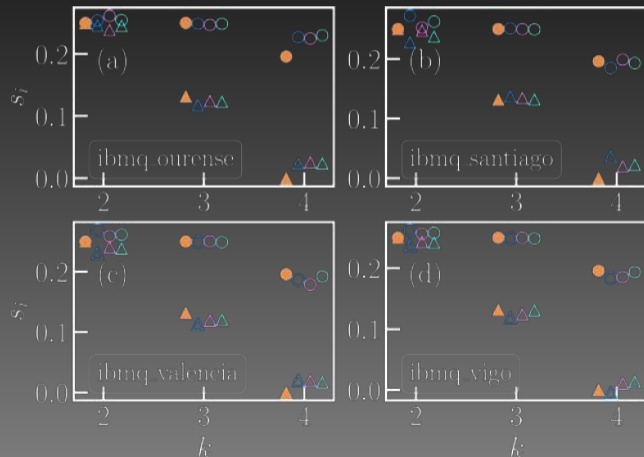
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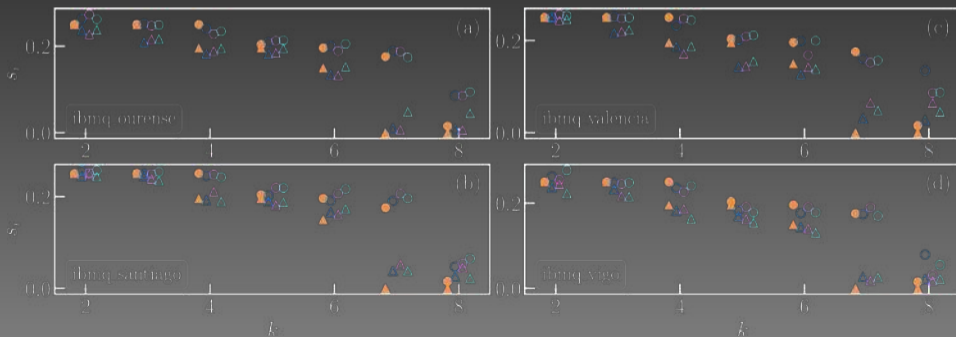
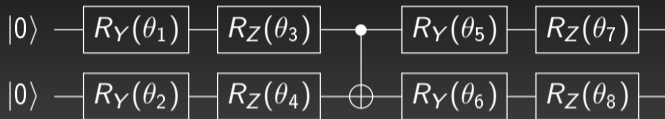
- Spectrum of S_k , $k \geq 2$



Dimensional Expressivity Analysis

Results for two qubits on quantum hardware

- Circuit we examine



Dimensional Expressivity Analysis

Summary

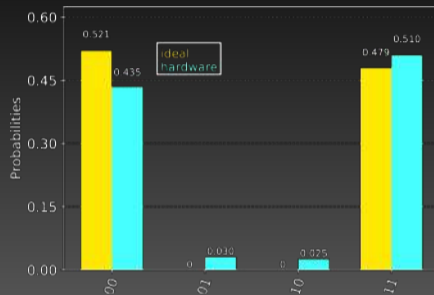
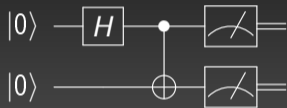
- Allows for optimizing a given circuit by identifying and **removing redundant parameters**
- Makes it possible to remove unwanted symmetries as well
- Can be **efficiently** performed using a **hybrid quantum-classical** approach

4.

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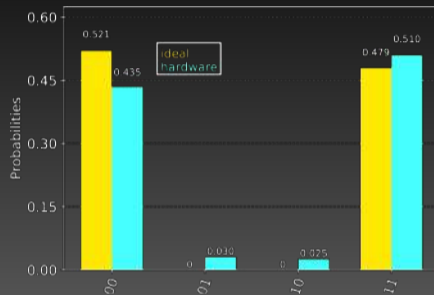
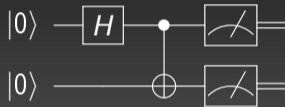
Measurement Error Mitigation

Noise on current quantum devices



Measurement Error Mitigation

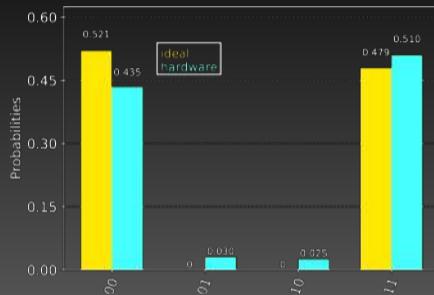
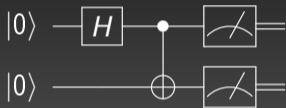
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- Errors arise from:
 - ▶ Imperfect gates and crosstalk
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 - ▶ Measurement/readout

Measurement Error Mitigation

Noise on current quantum devices



- Errors arise from:
 - ▶ Imperfect gates and crosstalk
 - ▶ Coupling to environment
 - ▶ Measurement/readout
- **Error mitigation:** Try to correct for (some of) these errors
- Typical error rates
 - ▶ Single-qubit gates: 0.1% - 0.3%
 - ▶ Two-qubit gates: 0.3% - 5%
 - ▶ Measurement/readout: 1% - 30%

Measurement Error Mitigation

Measurement Error Mitigation

- Focus on a simple **low-overhead**, **resource-efficient** mitigation scheme suitable even for small devices
- Assumptions:
 - ▶ **Only measurement errors**, no other sources of noise
 - ▶ **Uncorrelated bit flips**, readout errors are not correlated between qubits
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- Idea: construct random operators \tilde{O} such that such the expectation value subject to noise corresponds to the true expectation value

$$\mathbb{E}\langle\psi|\tilde{O}|\psi\rangle = \langle\psi|O|\psi\rangle$$

Measurement Error Mitigation

Single-qubit example

- Consider a single qubit with flip probabilities $0 \xrightarrow{p_0} 1$, $1 \xrightarrow{p_1} 0$ and measure the

$$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \text{ operator}$$

Readout	Bit Flips	Probability	Noisy operator
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Single-qubit example

- Consider a single qubit with flip probabilities $0 \xrightarrow{p_0} 1$, $1 \xrightarrow{p_1} 0$ and measure the

$$Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \text{ operator}$$

Readout	Bit Flips	Probability	Noisy operator
correct	$0 \rightarrow 0, 1 \rightarrow 1$	$(1 - p_0)(1 - p_1)$	$\tilde{Z} = Z$
incorrect	$0 \rightarrow 1, 1 \rightarrow 0$	$p_0 p_1$	$\tilde{Z} = -Z$
0 outcome incorrect	$0 \rightarrow 1, 1 \rightarrow 1$	$p_0(1 - p_1)$	$\tilde{Z} = -\mathbb{1}$
1 outcome incorrect	$0 \rightarrow 0, 1 \rightarrow 0$	$(1 - p_0)p_1$	$\tilde{Z} = \mathbb{1}$

- Expected value of the noisy operator

$$\begin{aligned} \mathbb{E}\tilde{Z} &= (1 - p_0)(1 - p_1)Z - p_0 p_1 Z - p_0(1 - p_1)\mathbb{1} + (1 - p_0)p_1\mathbb{1} \\ &= (1 - p_0 - p_1)Z + (p_0 - p_1)\mathbb{1} \end{aligned}$$

Measurement Error Mitigation

Single-qubit Example

- Expected value of the noisy operator

$$\mathbb{E}\tilde{Z} = (1 - p_0 - p_1)Z + (p_0 - p_1)\mathbb{1}$$

- Reconstruction of the true expectation value

available from calibration on noisy quantum device

$$\langle \psi | Z | \psi \rangle = \frac{1}{1 - p_0 - p_1} \times \mathbb{E} \langle \psi | \tilde{Z} | \psi \rangle - \frac{p_0 - p_1}{1 - p_0 - p_1}$$

true expectation value

measurement on noisy quantum device

Measurement Error Mitigation

Single-qubit Example

- Expected value of the noisy operator

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Diagram illustrating the reconstruction of the true expectation value from a noisy measurement. The equation shows the true expectation value $\langle \psi | Z | \psi \rangle$ (labeled "true expectation value") is equal to the measured expectation value $\mathbb{E} \langle \psi | \tilde{Z} | \psi \rangle$ (labeled "measurement on noisy quantum device") multiplied by the inverse of the probability of no flip, $\frac{1}{1 - p_0 - p_1}$, minus the bias term $\frac{p_0 - p_1}{1 - p_0 - p_1}$. Arrows indicate that the calibration probabilities p_0 and p_1 are available from calibration on a noisy quantum device.

- Calibration of the flip probabilities



Measurement Error Mitigation

Generalization to multiple qubits and arbitrary operators

- General formula for $\tilde{O}_k \in \{\tilde{\mathbb{1}}_k, \tilde{Z}_k\}$

$$\mathbb{E} \left(\tilde{O}_N \otimes \cdots \otimes \tilde{O}_1 \right) = \sum_{O \in \{\mathbb{1}, Z\}^{\otimes N}} \Gamma(O_N | \tilde{O}_N) O_N \otimes \cdots \otimes \Gamma(O_1 | \tilde{O}_1) O_1$$

where

$$\Gamma(O_q | \tilde{O}_q) = \begin{cases} 1 - p_{q,0} - p_{q,1} & \text{for } \tilde{O}_q = \tilde{Z}_q \wedge O_q = Z_q \\ p_{q,1} - p_{q,0} & \text{for } \tilde{O}_q = \tilde{Z}_q \wedge O_q = \mathbb{1}_q \\ 1 & \text{for } O_q = \mathbb{1}_q \wedge \tilde{O}_q = \tilde{\mathbb{1}}_q \\ 0 & \text{for } O_q = Z_q \wedge \tilde{O}_q = \tilde{\mathbb{1}}_q \end{cases}$$

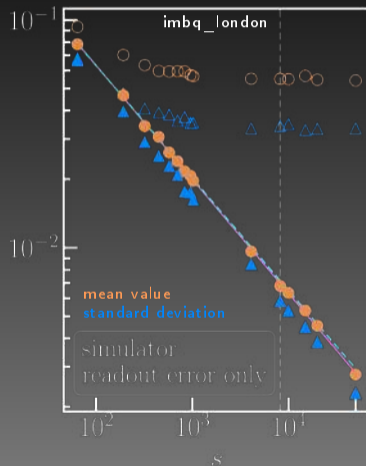
- Set of equations is **invertible** as long as $p_{q,0} + p_{q,1} \neq 1$

Measurement Error Mitigation

Two-qubit case, classical simulation

- Results for a **classical simulation** with **readout errors only**
- Measure the expectation value of $Z \otimes Z$ for 1050 random parameter sets
- Monitor the average and standard deviation of

$$|\langle \psi | Z \otimes Z | \psi \rangle_{\text{exact}} - \langle \psi | Z \otimes Z | \psi \rangle_{\text{mitigated}}|$$

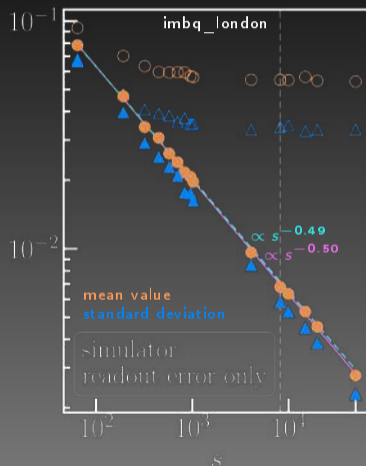


Measurement Error Mitigation

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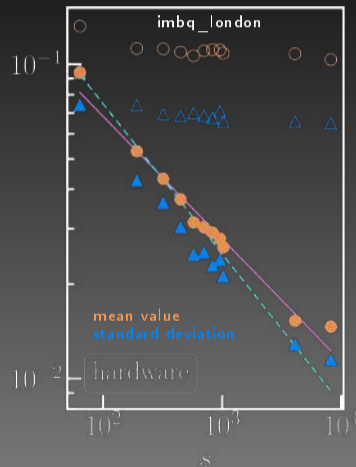
⇒ Mitigated results show power law decay $s^{-1/2}$ just as in the noise-free case

Measurement Error Mitigation

Two-qubit case, quantum hardware

- Results for **IBM quantum hardware**
- Measure the expectation value of $Z \otimes Z$ for 1050 random parameter sets
- Monitor the average and standard deviation of

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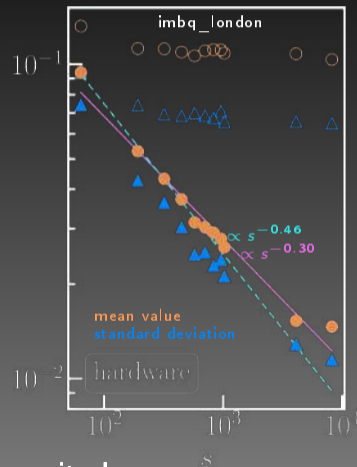


Measurement Error Mitigation

Two-qubit case, quantum hardware

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⇒ Improvement of the error up to one order of magnitude

Measurement Error Mitigation

Summary

- For **local Hamiltonians** the **overhead cost** is **polynomial**
- It is possible to do a probabilistic version of the mitigation scheme
- The idea of constructing random operators is very general and can potentially be applied to mitigate other kinds of errors
 - ▶ Incorporate correlations between qubits
 - ▶ Relaxation errors
 - ▶ ...

5.

- 1 Motivation
- 2 The Circuit Model of Quantum Computing
- 3 Dimensional Expressivity Analysis
- 4 Measurement Error Mitigation
- 5 Summary & Outlook

Summary & Outlook

Summary

- Noisy intermediate-scale quantum devices are available
- Hybrid-quantum classical algorithms are promising for these devices
- Dimensional Expressivity Analysis allows for designing minimal maximally expressive circuits for these applications
- Measurement/readout errors can be efficiently mitigated with low overhead

Summary & Outlook

Summary

- Noisy intermediate-scale quantum devices are available
- Hybrid-quantum classical algorithms are promising for these devices
- Dimensional Expressivity Analysis allows for designing minimal maximally expressive circuits for these applications
- Measurement/readout errors can be efficiently mitigated with low overhead

Outlook

- Generalize Dimensional Expressivity Analysis to be able to quantify the approximation error of a given ansatz
- Extend the mitigation scheme to various other types of error
- Quantum hardware is advancing quickly

Thank you for your attention!

Questions?

Appendix A: Dimensional Expressivity Analysis

Hybrid quantum-classical Dimensional Expressivity Analysis

- $\Re \langle \partial_j C(\vec{\theta}) | \partial_k C(\vec{\theta}) \rangle$ can be obtained on the quantum device

Appendix A: Dimensional Expressivity Analysis

Hybrid quantum-classical Dimensional Expressivity Analysis

- $\Re \langle \partial_j C(\vec{\theta}) | \partial_k C(\vec{\theta}) \rangle$ can be obtained on the quantum device
- In general $\Re \langle \psi | \phi \rangle$ can be measured using an ancilla qubit provided one can prepare the state

$$|\chi\rangle = \frac{1}{\sqrt{2}} (|0\rangle \otimes |\psi\rangle + |1\rangle \otimes |\phi\rangle)$$

- Applying a Hadamard gate on the ancilla one finds

$$(H \otimes \mathbf{1}) |\chi\rangle = \frac{1}{\sqrt{2}} (|0\rangle \otimes (|\psi\rangle + |\phi\rangle) + |1\rangle \otimes (|\psi\rangle - |\phi\rangle))$$

- Probability of measuring the ancilla in zero

$$p(\text{ancilla} = 0) = \frac{1}{2} (1 + \Re \langle \psi | \phi \rangle)$$